# Analysis of Optimization Algorithms via Sum-of-Squares

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# Motivation

Iterative algorithms using only (sub)gradient information

- Low computational complexity
- Ideal for large-scale problems with low-accuracy requirements

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 $\min_{\mathbf{x}\in\mathbb{R}^n}f(\mathbf{x})$ 

 $\mathsf{GD}: \mathbf{x}_{k+1} = \mathbf{x}_k - \gamma_k \nabla f(\mathbf{x}_k)$ 



Figure 1: Gradient Descent (GD)

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Performance  $Metric(\mathbf{x}_{k+1}) \leq t$  Performance  $Metric(\mathbf{x}_k)$ 

holds for all  $f \in \mathcal{F}$  and all  $\{\mathbf{x}_k\}_{k \geq 1}$  generated by  $\mathcal{A}$ 

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#### Performance metric

- Objective function accuracy  $f(\mathbf{x}_k) f(\mathbf{x}_*)$ , also denoted as  $f_k f_*$
- Squared distance to optimality  $\|\mathbf{x}_k \mathbf{x}_*\|^2$
- Squared residual gradient norm  $\|\nabla f(\mathbf{x}_k)\|^2$ , also denoted as  $\|\mathbf{g}_k\|^2$

#### L-smooth, $\mu$ -strongly convex (( $\mu$ , L)-smooth) function

• 
$$\|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \le L \|\mathbf{x} - \mathbf{y}\| \quad \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^n$$

• 
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**Theorem (de Klerk et al., 2017)** Given an  $(\mu, L)$ -smooth function, apply GD with exact line search:

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A sum-of-squares (SOS) proof

$$\left(\frac{L-\mu}{L+\mu}\right)^{2}(f_{k}-f_{*})-(f_{k+1}-f_{*}) \geq \frac{\mu}{4}\left(\frac{\|\mathbf{q}_{1}\|^{2}}{1+\sqrt{\mu/L}}+\frac{\|\mathbf{q}_{2}\|^{2}}{1-\sqrt{\mu/L}}\right) \geq 0$$

where  $\mathbf{q}_1$  and  $\mathbf{q}_2$  are linear functions of  $\mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, \mathbf{g}_k, \mathbf{g}_{k+1}$ 

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**GD** with constant step size,  $\gamma \in (0, \frac{2}{L})$  (Lessard et al., 2016)

$$\begin{split} &\rho^2 \|\mathbf{x}_k - \mathbf{x}_*\|^2 - \|\mathbf{x}_{k+1} - \mathbf{x}_*\|^2 \\ \geq \begin{cases} \frac{\gamma(2 - \gamma(L+\mu))}{L-\mu} \|\mathbf{g}_k - \mu(\mathbf{x}_k - \mathbf{x}_*)\|^2, & \gamma \in \left(0, \frac{2}{L+\mu}\right] \\ \frac{\gamma(\gamma(L+\mu) - 2)}{L-\mu} \|\mathbf{g}_k - L(\mathbf{x}_k - \mathbf{x}_*)\|^2, & \gamma \in \left[\frac{2}{L+\mu}, \frac{2}{L}\right) \end{cases} \end{split}$$

where  $\rho:=\max\{\left|1-\gamma\mu\right|,\left|1-\gamma L\right|\}$ 

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where  $\rho := \max\{\left|1 - \gamma \mu\right|, \left|1 - \gamma L\right|\}$ 

#### GD with the Armijo rule

$$\left(1-\frac{4\mu\epsilon(1-\epsilon)}{\eta L}\right)(f_k-f_*)-(f_{k+1}-f_*)\geq \frac{2\epsilon(1-\epsilon)}{\eta(L-\mu)}\|\mathbf{g}_k+\mu(\mathbf{x}_*-\mathbf{x}_k)\|^2,$$

where  $\epsilon \in (0,1), \eta > 1$  are the parameters of the Armijo rule

### More SOS proofs: Proximal Gradient Method (PGM)

PGM with exact line search (Taylor et al., 2018)

$$\left(\frac{L-\mu}{L+\mu}\right)^2 (f_k - f_*) - (f_{k+1} - f_*) \ge \begin{pmatrix} 1 \\ z \end{pmatrix}^\top Q^* \begin{pmatrix} 1 \\ z \end{pmatrix}$$

where  $\mathbf{z} = (f_*, f_k, f_{k+1}, \mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, \dots)$  and  $Q^*$  is a PSD matrix

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where  $\mathbf{z} = (f_*, f_k, f_{k+1}, \mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, ...)$  and  $Q^*$  is a PSD matrix **PGM with constant step size,**  $\gamma \in (0, \frac{2}{L})$  (Taylor et al., 2018)

$$\begin{split} &\rho^{2} \|\mathbf{r}_{k} + \mathbf{s}_{k}\|^{2} - \|\mathbf{r}_{k+1} + \bar{\mathbf{s}}_{k+1}\|^{2} \\ &\geq \begin{cases} \rho^{2} \|\mathbf{q}_{1}\|^{2} + \frac{2 - \gamma(L+\mu)}{\gamma(L-\mu)} \|\mathbf{q}_{2}\|^{2}, & \gamma \in \left(0, \frac{2}{L+\mu}\right) \\ \rho^{2} \|\mathbf{q}_{1}\|^{2} + \frac{\gamma(L+\mu) - 2}{\gamma(L-\mu)} \|\mathbf{q}_{3}\|^{2}, & \gamma \in \left[\frac{2}{L+\mu}, \frac{2}{L}\right) \end{cases} \end{split}$$

where  $\rho := \max\{|1 - \gamma \mu|, |1 - \gamma L|\}$  and  $\mathbf{q}_1$ ,  $\mathbf{q}_2$  and  $\mathbf{q}_3$  are linear functions of  $\mathbf{s}_k, \mathbf{\bar{s}}_{k+1}, \mathbf{r}_k, \mathbf{r}_{k+1}$ 

# **Related Work**

$$\max f_{N} - f_{*}$$
  
s.t.  $f \in \mathcal{F}$   
 $\mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_{k}, f_{k}, \mathbf{g}_{k}), \ k = 0, \dots, N-1$   
 $f_{0} - f_{*} \leq R$ 

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- GD with exact line search (de Klerk et al., 2017)
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- Proximal Point Algorithm, Conditional Gradient Method, ... (Taylor et al., 2017a); ...
- Formulated as an SDP (exact under certain conditions)
- Inspiration for our current work



 $\begin{array}{lll} \text{State:} & \xi_{k+1} = A\xi_k + Bu_k \\ \text{Output:} & y_k = C\xi_k + Du_k \\ \text{Feedback:} & u_k = \phi(y_k) \end{array}$ 



• For first-order methods,  $u_k = \phi(y_k) := \nabla f(y_k)$ 



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- Convergence of algorithm  $\equiv$  stability of dynamical system
- Replace  $\phi$  with quadratic constraints on all instances of  $(y_k, u_k)$

# The Sum-of-Squares Approach

**Sum-of-squares of degree-2 polynomials** For z = (x, y) and d = 2,

$$p(\mathbf{z}) = (x^2 - 2)^2 + (x - 3y)^2$$
$$= x^4 - 3x^2 - 6xy + 9y^2 + 4$$

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=  $\begin{pmatrix} 1 \\ y \\ x^2 \\ xy \\ y^2 \end{pmatrix}^\top \begin{pmatrix} Q_{1,1} & Q_{1,2} & \dots \\ Q_{2,1} & \ddots \\ \vdots \end{pmatrix} \begin{pmatrix} 1 \\ y \\ x^2 \\ xy \\ y^2 \end{pmatrix}$ 

where  $Q \succeq 0$ 

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$$1: \quad 4 = Q_{1,1}$$
  
x: 
$$0 = Q_{1,2} + Q_{2,1} = 2Q_{1,2}$$

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where  $Q \succeq 0$ 

- Match coefficients of  $1, x, y, x^2, \dots, y^4 \implies$  affine constraints
- Affine constraints +  $Q \succeq 0$  condition  $\implies$  SDP feasibility problem

min 0 s.t.  $Q \succeq 0$ , affine constraints

$$K = \{ \mathbf{z} : h_i(\mathbf{z}) \ge 0, v_j(\mathbf{z}) = 0 \quad \forall i, j \}$$

for some polynomials  $h_i(\mathbf{z}), v_j(\mathbf{z})$ 's

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#### "Constrained" nonnegativity

To certify that  $p(\mathbf{z}) \ge 0$  for all  $\mathbf{z} \in K$  (Parrilo, 2003; Lasserre, 2007),

$$p(\mathbf{z}) = \sigma_0(\mathbf{z}) + \sum_i \sigma_i(\mathbf{z}) h_i(\mathbf{z}) + \sum_j \theta_j(\mathbf{z}) v_j(\mathbf{z})$$

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where a sum-of-squares (SOS) polynomial:  $\sigma(\mathbf{z}) = \sum_{k} q_{k}^{2}(\mathbf{z})$ 

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- Fix degrees of the  $\sigma_i(\mathbf{z})$ 's and  $\theta_j(\mathbf{z})$ 's
- Construct and solve the corresponding SDP feasibility problem
# Analyzing optimization algorithms

$$\begin{array}{ll} \min \ t \\ \text{s.t.} \ t(f_k - f_*) - (f_{k+1} - f_*) \ge 0 \\ t \in (0, 1) \\ f \in \mathcal{F} \\ \mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_k, f_k, \mathbf{g}_k) \end{array}$$

where  $f_k = f(\mathbf{x}_k)$  and  $\mathbf{g}_k \in \partial f(\mathbf{x}_k)$  for all k

min t  
s.t. 
$$t(f_k - f_*) - (f_{k+1} - f_*) \ge 0$$
  
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 $h_i(\mathbf{z}) \ge 0, \ \forall i$   
 $v_i(\mathbf{z}) = 0, \ \forall j$ 

where  $\mathbf{z} = (f_*, f_k, f_{k+1}, \mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, \mathbf{g}_*, \mathbf{g}_k, \mathbf{g}_{k+1}) \in \mathbb{R}^{6n+3}$ 

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- Identify polynomial inequalities h<sub>i</sub>(z) ≥ 0 and equalities v<sub>j</sub>(z) = 0 necessarily satisfied given F and A
- E.g. f is L-smooth  $\implies \|\nabla f(\mathbf{x}) \nabla f(\mathbf{y})\| \le L \|\mathbf{x} \mathbf{y}\| \quad \forall \mathbf{x}, \mathbf{y}$

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E.g. f is L-smooth 
$$\implies \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \le L \|\mathbf{x} - \mathbf{y}\| \quad \forall \mathbf{x}, \mathbf{y}$$
  
 $\implies h_1(\mathbf{z}) = L^2 \|\mathbf{x}_k - \mathbf{x}_*\|^2 - \|\mathbf{g}_k - \mathbf{g}_*\|^2 \ge 0$ 

$$\begin{array}{l} \min t \\ \text{s.t.} \quad t(f_k - f_*) - (f_{k+1} - f_*) \ge 0 \\ \quad t \in (0, 1) \\ \quad h_i(\mathbf{z}) \ge 0, \ \forall i \\ \quad \mathbf{v}_i(\mathbf{z}) = 0, \ \forall j \end{array} \right\} K$$

where  $\mathbf{z} = (f_*, f_k, f_{k+1}, \mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, \mathbf{g}_*, \mathbf{g}_k, \mathbf{g}_{k+1}) \in \mathbb{R}^{6n+3}$ 

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where  $K = \{ \mathbf{z} : h_i(\mathbf{z}) \ge 0, v_j(\mathbf{z}) = 0 \quad \forall i, j \}$ 

### Second relaxation step

 $\begin{array}{ll} \min & t \\ \text{s.t.} & t(f_k - f_*) - (f_{k+1} - f_*) \geq 0 \quad \forall \; \mathbf{z} \in K \\ & t \in (0, 1) \end{array}$ 

where  $K = \{ \mathbf{z} : h_i(\mathbf{z}) \ge 0, v_j(\mathbf{z}) = 0 \quad \forall i, j \}$ 

2. Relax the constraint that  $p(\mathbf{z}) := t(f_k - f_*) - (f_{k+1} - f_*)$  is nonnegative over K with an SOS cert

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2. Relax nonnegativity constraint with an SOS cert, i.e., relax

$$p(\mathbf{z}) \geq 0 \,\,\forall \,\, \mathbf{z} \in K \,\, ext{to} \,\, p(\mathbf{z}) = \sigma_0(\mathbf{z}) + \sum_i \sigma_i(\mathbf{z}) h_i(\mathbf{z}) + \sum_j \theta_j(\mathbf{z}) v_j(\mathbf{z})$$

for SOS polynomials  $\sigma_i(\mathbf{z})$  and arbitrary polynomials  $\theta_j(\mathbf{z})$ 

1. Identify polynomial (in)equalities necessarily satisfied given  ${\cal F}$  and  ${\cal A}$ 

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ight\}$$

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4. Generalize this to a feasible SDP solution for the multivariate case

# Example

# **Function class**

Class of  $(\mu, L)$ -smooth functions (denoted  $\mathcal{F}_{\mu,L}(\mathbb{R}^n)$ )

### **Function class**

Class of  $(\mu, L)$ -smooth functions (denoted  $\mathcal{F}_{\mu,L}(\mathbb{R}^n)$ ) **Theorem (\mathcal{F}\_{\mu,L}-interpolability (Taylor et al., 2017b))**  *Given a set* { $(\mathbf{x}_i, f_i, \mathbf{g}_i)$ }<sub> $i \in I$ </sub>,  $\exists$  *a* ( $\mu, L$ )-smooth function *f* where  $f_i = f(\mathbf{x}_i)$  and  $\mathbf{g}_i \in \partial f(\mathbf{x}_i) \forall i$  iff

$$f_i - f_j - \mathbf{g}_j^{\top}(\mathbf{x}_i - \mathbf{x}_j) \ge \frac{L}{2(L - \mu)} \left[ \frac{1}{L} \|\mathbf{g}_i - \mathbf{g}_j\|^2 + \mu \|\mathbf{x}_i - \mathbf{x}_j\|^2 - 2\frac{\mu}{L} (\mathbf{g}_j - \mathbf{g}_i)^{\top} (\mathbf{x}_j - \mathbf{x}_i) \right] \forall i \neq j$$

### **Function class**

Class of  $(\mu, L)$ -smooth functions (denoted  $\mathcal{F}_{\mu,L}(\mathbb{R}^n)$ )

**Theorem** ( $\mathcal{F}_{\mu,L}$ -interpolability (Taylor et al., 2017b)) Given a set {( $\mathbf{x}_i, f_i, \mathbf{g}_i$ )}<sub>i \in I</sub>,  $\exists$  a ( $\mu, L$ )-smooth function f where  $f_i = f(\mathbf{x}_i)$  and  $\mathbf{g}_i \in \partial f(\mathbf{x}_i) \forall i$  iff

$$f_{i} - f_{j} - \mathbf{g}_{j}^{\top}(\mathbf{x}_{i} - \mathbf{x}_{j}) \geq \frac{L}{2(L-\mu)} \left[ \frac{1}{L} \|\mathbf{g}_{i} - \mathbf{g}_{j}\|^{2} + \mu \|\mathbf{x}_{i} - \mathbf{x}_{j}\|^{2} - 2\frac{\mu}{L}(\mathbf{g}_{j} - \mathbf{g}_{i})^{\top}(\mathbf{x}_{j} - \mathbf{x}_{i}) \right] \forall i \neq j$$

E.g. if we set i = k and j = k + 1:

$$\begin{aligned} f_{k} - f_{k+1} - \mathbf{g}_{k+1}^{\top}(\mathbf{x}_{k} - \mathbf{x}_{k+1}) - \frac{L}{2(L-\mu)} \bigg[ \frac{1}{L} \|\mathbf{g}_{k} - \mathbf{g}_{k+1}\|^{2} + \mu \|\mathbf{x}_{k} - \mathbf{x}_{k+1}\|^{2} \\ &- 2\frac{\mu}{L} (\mathbf{g}_{k+1} - \mathbf{g}_{k})^{\top} (\mathbf{x}_{k+1} - \mathbf{x}_{k}) \bigg] \geq 0 \end{aligned}$$

### Gradient Descent (GD) with exact line search (ELS)

$$\gamma_k = \arg\min_{\gamma>0} f\left(\mathbf{x}_k - \gamma \mathbf{g}_k\right)$$

$$\mathbf{x}_{k+1} = \mathbf{x}_k - \gamma_k \mathbf{g}_k$$

where  $\mathbf{g}_k = \nabla f(\mathbf{x}_k)$ 

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**Derived constraints** 

$$egin{aligned} \mathbf{g}_{k+1}^{ op}\left(\mathbf{x}_{k+1}-\mathbf{x}_{k}
ight) &= 0 \ \mathbf{g}_{k+1}^{ op}\mathbf{g}_{k} &= 0 \end{aligned}$$

- 6 *F*<sub>µ,L</sub>-interpolability constraints corresponding to k, k + 1, \* (h<sub>i</sub>(z) ≥ 0)
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### The SDP

s.t. 
$$p(\mathbf{z}) = \sigma_0(\mathbf{z}) + \sum_i \sigma_i(\mathbf{z})h_i(\mathbf{z}) + \sum_j \theta_j(\mathbf{z})v_j(\mathbf{z})$$
  
 $t \in (0, 1)$   
 $\sigma_0(\mathbf{z})$ : SOS of linear polynomials  
 $\sigma_i(\mathbf{z}), i > 1$ : SOS of degree 0,  $\theta_i(\mathbf{z})$ : degree 0

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s.t. 
$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}^{\top} Q \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix} + \sum_{i} \sigma_{i}(\mathbf{z}) h_{i}(\mathbf{z}) + \sum_{j} \theta_{j}(\mathbf{z}) v_{j}(\mathbf{z})$$
  
 $t \in (0, 1)$   
 $Q \succeq 0$   
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 $t \in (0, 1)$   
 $Q \succeq 0$   
 $\sigma_{i} \geq 0, i \geq 1, \ \theta_{i} \in \mathbb{R}$ 

$$t(f_{k} - f_{*}) - (f_{k+1} - f_{*})$$

$$= \begin{pmatrix} 1 \\ f_{*} \\ \vdots \end{pmatrix}^{\top} \begin{pmatrix} Q_{1,1} & Q_{1,2} & \cdots \\ Q_{2,1} & \ddots & \\ \vdots & \cdots & \end{pmatrix} \begin{pmatrix} 1 \\ f_{*} \\ \vdots \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j} v_{j}(\mathbf{z})$$

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$$1 - t =$$

$$t(f_{k} - f_{*}) - (f_{k+1} - f_{*})$$

$$= \begin{pmatrix} \mathbf{1} \\ f_{*} \\ \vdots \end{pmatrix}^{\top} \begin{pmatrix} Q_{1,1} & Q_{1,2} & \cdots \\ Q_{2,1} & \ddots & \\ \vdots & & \end{pmatrix} \begin{pmatrix} \mathbf{1} \\ f_{*} \\ \vdots \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j} v_{j}(\mathbf{z})$$

$$1 - t = 2Q_{1,2}$$

$$t(f_{k} - f_{*}) - (f_{k+1} - f_{*})$$

$$= \begin{pmatrix} 1 \\ f_{*} \\ \vdots \end{pmatrix}^{\top} \begin{pmatrix} Q_{1,1} & Q_{1,2} & \cdots \\ Q_{2,1} & \ddots & \\ \vdots & \cdots \end{pmatrix} \begin{pmatrix} 1 \\ f_{*} \\ \vdots \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j} v_{j}(\mathbf{z})$$

Matching coefficients for e.g.  $f_*$  term:

$$1 - t = 2Q_{1,2} - \sigma_2$$

Constraint  $h_2(\mathbf{z}) \geq 0$ :

$$f_k - f_* - \frac{L}{2(L-\mu)} \left[ \frac{1}{L} \|\mathbf{g}_k\|^2 + \mu \|\mathbf{x}_k - \mathbf{x}_*\|^2 + 2\frac{\mu}{L} \mathbf{g}_k^\top (\mathbf{x}_* - \mathbf{x}_k) \right] \ge 0$$

$$t(f_k - f_*) - (f_{k+1} - f_*)$$

$$= \begin{pmatrix} 1 \\ f_* \\ \vdots \end{pmatrix}^\top \begin{pmatrix} Q_{1,1} & Q_{1,2} & \cdots \\ Q_{2,1} & \ddots & \\ \vdots & \cdots & \end{pmatrix} \begin{pmatrix} 1 \\ f_* \\ \vdots \end{pmatrix} + \sum_i \sigma_i h_i(\mathbf{z}) + \sum_j \theta_j v_j(\mathbf{z})$$

$$1 - t = 2Q_{1,2} - \sigma_2 - \sigma_4 + \sigma_5 + \sigma_6$$

$$t(f_{k} - f_{*}) - (f_{k+1} - f_{*}) = \begin{pmatrix} 1 \\ z \end{pmatrix}^{\top} Q \begin{pmatrix} 1 \\ z \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(z) + \sum_{j} \theta_{j} v_{j}(z)$$

$$1: \quad 0 = Q_{1,1}$$

$$f_{*}: \quad 1 - t = 2Q_{1,2} - \sigma_{2} - \sigma_{4} + \sigma_{5} + \sigma_{6}$$

$$f_{k}: \quad t = 2Q_{1,3} + \sigma_{1} + \sigma_{2} - \sigma_{3} - \sigma_{5}$$

$$f_{k+1}: \quad -1 = 2Q_{1,4} - \sigma_{1} + \sigma_{3} + \sigma_{4} - \sigma_{6}$$

$$:$$

# Solving the SDP

s.t. 
$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}^{\top} Q \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j} v_{j}(\mathbf{z})$$
  
 $t \in (0, 1)$   
 $Q \succeq 0$   
 $\sigma_{i} \geq 0, \ \theta_{i} \in \mathbb{R}$
# Solving the SDP

min t

s.t. 
$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}^{\top} Q \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix} + \sum_{i} \sigma_{i} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j} v_{j}(\mathbf{z})$$
  
 $t \in (0, 1)$   
 $Q \succeq 0$   
 $\sigma_{i} > 0, \ \theta_{i} \in \mathbb{R}$ 

Guessing the optimal solution

$$\begin{split} t^* &= \left(\frac{L-\mu}{L+\mu}\right)^2, \ Q^* = \left(\frac{0_{4\times 4}}{0_{5\times 4}} \left|\frac{0_{4\times 5}}{0_{5\times 5}}\right), \ \sigma_1^* = \frac{L-\mu}{L+\mu}, \ \dots \\ Q^*_{5\times 5} &= \left(\begin{array}{ccc} \frac{2L^2\mu^2}{(L+\mu)^2(L-\mu)} & -\frac{L\mu^2}{(L+\mu)^2} & -\frac{L\mu^2}{(L+\mu)^2} & \frac{L\mu}{(L+\mu)(L-\mu)} \\ & \frac{L\mu(L+3\mu)}{2(L+\mu)^2} & -\frac{L\mu}{2(L+\mu)} & -\frac{\mu}{2(L+\mu)} \\ & \frac{L\mu}{2(L-\mu)} & \frac{\mu}{2(L+\mu)} & -\frac{\mu}{2(L-\mu)} \\ & \frac{L+3\mu}{2(L-\mu)} & \frac{L+3\mu}{2(L+\mu)^2} & \frac{1}{2(L+\mu)} \\ & \frac{L+3\mu}{2(L-\mu)} & \frac{1}{2(L-\mu)} \end{array} \right) \end{split}$$

## Multivariate case

Recall 
$$z = (f_*, f_k, f_{k+1}, x_*, x_k, x_{k+1}, g_*, g_k, g_{k+1})$$
  
Let  $z = (z_0, z_1, \dots, z_n)$  where

 $\textbf{z}_0 = (f_*, f_k, f_{k+1}) \text{ and } \textbf{z}_\ell = (x_*(\ell), x_k(\ell), x_{k+1}(\ell), g_*(\ell), g_k(\ell), g_{k+1}(\ell))$ 

## Multivariate case

Recall 
$$\mathbf{z} = (f_*, f_k, f_{k+1}, \mathbf{x}_*, \mathbf{x}_k, \mathbf{x}_{k+1}, \mathbf{g}_*, \mathbf{g}_k, \mathbf{g}_{k+1})$$
  
Let  $\mathbf{z} = (\mathbf{z}_0, \mathbf{z}_1, \dots, \mathbf{z}_n)$  where

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Require polynomials  $p(\mathbf{z})$ ,  $h_i(\mathbf{z})$  and  $v_j(\mathbf{z})$  satisfy

$$\mathsf{pol}(z) = \mathsf{pol}^0(z_0) + \sum_{\ell=1}^n \mathsf{pol}^1(z_\ell)$$

Recall 
$$z = (f_*, f_k, f_{k+1}, x_*, x_k, x_{k+1}, g_*, g_k, g_{k+1})$$
  
Let  $z = (z_0, z_1, \dots, z_n)$  where

 $\mathbf{z}_0 = (f_*, f_k, f_{k+1}) \text{ and } \mathbf{z}_\ell = (x_*(\ell), x_k(\ell), x_{k+1}(\ell), g_*(\ell), g_k(\ell), g_{k+1}(\ell))$ 

Require polynomials  $p(\mathbf{z})$ ,  $h_i(\mathbf{z})$  and  $v_j(\mathbf{z})$  satisfy

$$\mathsf{pol}(\mathsf{z}) = \mathsf{pol}^0(\mathsf{z}_0) + \sum_{\ell=1}^n \mathsf{pol}^1(\mathsf{z}_\ell)$$

• Fulfilled when polynomials are linear in *f*'s and in the inner products of **x**'s and **g**'s e.g.

$$f_k - f_* - \frac{L}{2(L-\mu)} \left[ \frac{1}{L} \|\mathbf{g}_k\|^2 + \mu \|\mathbf{x}_k - \mathbf{x}_*\|^2 + 2\frac{\mu}{L} \mathbf{g}_k^\top (\mathbf{x}_{k+1} - \mathbf{x}_k) \right] \ge 0$$

#### Univariate case

$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \end{pmatrix}^\top \begin{pmatrix} 0 & & \\ & Q_0 & \\ & & Q_1 \end{pmatrix} \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \end{pmatrix} + \sum_i \sigma_i h_i(\mathbf{z}) + \sum_j \theta_j v_j(\mathbf{z})$$

#### Univariate case

$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \end{pmatrix}^\top \begin{pmatrix} 0 & & \\ & Q_0 & \\ & & Q_1 \end{pmatrix} \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \end{pmatrix} + \sum_i \sigma_i h_i(\mathbf{z}) + \sum_j \theta_j v_j(\mathbf{z})$$

#### Multivariate case

$$p(\mathbf{z}) = \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \\ \vdots \\ \mathbf{z}_n \end{pmatrix}^\top \begin{pmatrix} 0 & & \\ & Q_0 & \\ & & Q_1 \otimes I_n \end{pmatrix} \begin{pmatrix} 1 \\ \mathbf{z}_0 \\ \mathbf{z}_1 \\ \vdots \\ \mathbf{z}_n \end{pmatrix} + \sum_i \sigma_i h_i(\mathbf{z}) + \sum_j \theta_j v_j(\mathbf{z})$$

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#### Multivariate case

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SOS cert for univariate case  $\implies$  SOS cert for multivariate case

# Obtaining the SOS certificate

$$t^*(f_k - f_*) - (f_{k+1} - f_*) \stackrel{(1)}{=} \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix}^\top Q^* \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix} + \sum_i \sigma_i^* h_i(\mathsf{z}) + \sum_j \theta_j^* v_j(\mathsf{z})$$

1. Verifying the solution to be feasible

# Obtaining the SOS certificate

$$t^{*}(f_{k} - f_{*}) - (f_{k+1} - f_{*}) \stackrel{(1)}{=} \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix}^{\top} Q^{*} \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix} + \sum_{i} \sigma_{i}^{*} h_{i}(\mathsf{z}) + \sum_{j} \theta_{j}^{*} v_{j}(\mathsf{z})$$
$$\stackrel{(2)}{\geq} \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix}^{\top} Q^{*} \begin{pmatrix} 1 \\ \mathsf{z} \end{pmatrix}$$

- 1. Verifying the solution to be feasible
- 2. Since  $\sigma_i^*$ ,  $h_i(\mathbf{z}) \ge 0$  and  $v_j(\mathbf{z}) = 0$

# Obtaining the SOS certificate

$$t^{*}(f_{k} - f_{*}) - (f_{k+1} - f_{*}) \stackrel{(1)}{=} \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}^{\top} Q^{*} \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix} + \sum_{i} \sigma_{i}^{*} h_{i}(\mathbf{z}) + \sum_{j} \theta_{j}^{*} v_{j}(\mathbf{z})$$
$$\stackrel{(2)}{\geq} \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}^{\top} Q^{*} \begin{pmatrix} 1 \\ \mathbf{z} \end{pmatrix}$$
$$\stackrel{(3)}{=} \frac{\mu}{4} \left( \frac{\|\mathbf{q}_{1}\|^{2}}{1 + \sqrt{\mu/L}} + \frac{\|\mathbf{q}_{2}\|^{2}}{1 - \sqrt{\mu/L}} \right)$$

- 1. Verifying the solution to be feasible
- 2. Since  $\sigma_i^*$ ,  $h_i(\mathbf{z}) \ge 0$  and  $v_j(\mathbf{z}) = 0$
- 3. Expanding the quadratic term into its SOS decomposition

# Revisiting the PEP

$$\max f_{N} - f_{*}$$
  
s.t.  $f \in \mathcal{F}$   
 $\mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_{k}, f_{k}, \mathbf{g}_{k}), \ k = 0, \dots, N-1$   
 $f_{0} - f_{*} \leq R$ 

$$\max f_N - f_*$$
  
s.t.  $f \in \mathcal{F}$   
 $\mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_k, f_k, \mathbf{g}_k), \ k = 0, \dots, N-1$   
 $f_0 - f_* < R$ 

• Assume black-box model using oracle  $\mathcal{O}_f$ 

$$\max f_{N} - f_{*}$$
  
s.t.  $\exists f \in \mathcal{F}$  s.t.  $\mathcal{O}_{f}(\mathbf{x}_{i}) = \{f_{i}, \mathbf{g}_{i}\} \forall i$   
 $\mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_{k}, f_{k}, \mathbf{g}_{k}), \ k = 0, \dots, N-1$   
 $f_{0} - f_{*} \leq R$ 

• Assume black-box model using oracle  $\mathcal{O}_f$ 

$$\max f_{N} - f_{*}$$
  
s.t.  $\exists f \in \mathcal{F}$  s.t.  $\mathcal{O}_{f}(\mathbf{x}_{i}) = \{f_{i}, \mathbf{g}_{i}\} \forall i$   
 $\mathbf{x}_{k+1} = \mathcal{A}(\mathbf{x}_{k}, f_{k}, \mathbf{g}_{k}), \ k = 0, \dots, N-1$   
 $f_{0} - f_{*} \leq R$ 

- Assume black-box model using oracle  $\mathcal{O}_f$
- Identify polynomial inequalities h<sub>i</sub>(z) ≥ 0 and equalities v<sub>j</sub>(z) = 0 necessarily satisfied given F and A

 $\max f_N - f_*$ s.t.  $h_i(\mathbf{z}) \ge 0, \forall i$  $v_j(\mathbf{z}) = 0, \forall j$  $f_0 - f_* \le R$ 

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- Formulated as an SDP

 $\begin{array}{l} \max \ f_{\mathcal{N}} - f_{*} \\ \text{s.t.} \ \left\{ \mathbf{x}_{i} \right\} \text{ generated by } \mathcal{A} \\ f \in \mathcal{F} \\ f_{0} - f_{*} \leq R \\ \end{array} \tag{PEP}$ 

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N-step PEP-SDP



N-step PEP-SDP

↓ 1-step PEP-SDP  $\begin{array}{ccc} \max f_{N} - f_{*} & \min t \\ \text{s.t. } \{\mathbf{x}_{i}\} \text{ generated by } \mathcal{A} & \text{ f} \in \mathcal{F} \\ f_{0} - f_{*} \leq R & f \in \mathcal{F} \\ \end{array} & \begin{array}{c} \min t \\ \text{s.t. } t(f_{k} - f_{*}) - (f_{k+1} - f_{*}) \geq 0, \\ \forall \{\mathbf{x}_{i}\} \text{ generated by } \mathcal{A} \text{ and} \\ f \in \mathcal{F} & f \in \mathcal{F} \\ \end{array} \\ (\text{PEP}) & (\text{SOS}) \end{array}$ 

↓ N-step PEP-SDP ↓ 1-step PEP-SDP



$\begin{array}{l} \max \ f_N - f_* \\ \text{s.t.} \ \{ \mathbf{x}_i \} \text{ generated by } \mathcal{A} \\ f \in \mathcal{F} \\ f_0 - f_* \leq R \end{array} \\ (PEP) \end{array}$	$\begin{array}{c} \mbox{min } t \\ \mbox{s.t. } t(f_k - f_*) - (f_{k+1} - f_*) \ge 0, \\ & \forall \ \{ {\bf x}_i \} \ \mbox{generated by } \mathcal{A} \ \mbox{and} \\ & f \in \mathcal{F} \end{array}$ $(SOS)$
↓	↓
N-step PEP-SDP	Deg-d SOS-SDP
$\downarrow$	$\downarrow$
1-step PEP-SDP	Deg-1 SOS-SDP

$\begin{array}{l} \max \ f_N - f_* \\ \text{s.t.} \ \{ {\bf x}_i \} \text{ generated by } \mathcal{A} \\ f \in \mathcal{F} \\ f_0 - f_* \leq R \\ \end{array} \tag{PEP}$	$\begin{array}{l} \mbox{min } t \\ \mbox{s.t. } t(f_k - f_*) - (f_{k+1} - f_*) \geq 0, \\ & \forall \; \{ {\bf x}_i \} \; \mbox{generated by } \mathcal{A} \; \mbox{and} \\ & f \in \mathcal{F} \end{array} \\ \end{tabular} \tag{SOS}$
$\downarrow$	$\downarrow$
N-step PEP-SDP	Deg-d SOS-SDP
$\downarrow$	$\downarrow$
1-step PEP-SDP 🛛 🚝	$Dual \Longrightarrow Deg-1 SOS-SDP$

# Results

Form of bounds: Performance metric( $\mathbf{x}_{k+1}$ )  $\leq t^*$  Performance metric( $\mathbf{x}_k$ ) Set  $\epsilon$  and  $\eta$  to be parameters related to the Armijo rule, and

$$\rho_{\gamma} := \max\{\left|1 - \gamma \mu\right|, \left|1 - \gamma L\right|\}$$

Method	Step size $\gamma$	Metric	Rate t*	Known?
GD	ELS	$f_k - f_*$	$\left(\frac{L-\mu}{L+\mu}\right)^2$	de Klerk et al. (2017)

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	$\gamma \in (0, \frac{2}{L})$	$\ \mathbf{x}_k - \mathbf{x}_*\ ^2$	$ ho_{\gamma}^2$	Lessard et al. (2016)
	Armijo rule	$f_k - f_*$	$1 - rac{4\mu\epsilon(1-\epsilon)}{\eta L}$	New

## GD with Armijo-terminated Line Search

Our result  $t_{new}$  vs. known contraction factors  $t_{LY}$  (Luenberger and Ye, 2016) and  $t_{nemi}$  (Nemirovski, 1999):

$$t_{\mathsf{new}} = 1 - \frac{4\epsilon(1-\epsilon)}{\eta\kappa} \quad t_{\mathsf{LY}} = 1 - \frac{2\epsilon}{\eta\kappa} \quad t_{\mathsf{nemi}} = \frac{\kappa - (2-\epsilon^{-1})(1-\epsilon)\eta^{-1}}{\kappa + (\epsilon^{-1}-1)\eta^{-1}}$$

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**Figure 2:**  $t_{new}$  and  $t_{LY}$  for  $\epsilon = 0.25$ 

Condition number,  $\kappa = L/\mu$ 

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**Figure 2:**  $t_{new}$  and  $t_{LY}$  for  $\epsilon = 0.25$ 

**Figure 3:**  $t_{\text{new}}$  and  $t_{\text{nemi}}$  for  $\epsilon = 0.5$ 

Condition number,  $\kappa = L/\mu$ 

#### Composite convex minimization problem

$$\min_{\mathbf{x}\in\mathbb{R}^n}\left\{f(\mathbf{x}):=a(\mathbf{x})+b(\mathbf{x})\right\},\,$$

where  $a \in \mathcal{F}_{\mu,L}$  and b is closed, convex and proper.

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#### **Further assumptions**

Assume that the proximal operator of b,

$$\operatorname{prox}_{\gamma b}(\mathbf{x}) := \arg\min_{\mathbf{y}\in\mathbb{R}^n} \left\{ \gamma b(\mathbf{y}) + \frac{1}{2} \|\mathbf{x} - \mathbf{y}\|^2 \right\},$$

exists at every x

#### Proximal gradient method (PGM) with constant step size $\gamma > 0$

$$\mathbf{x}_{k+1} = \mathbf{prox}_{\gamma b} \left( \mathbf{x}_k - \gamma \nabla \mathbf{a}(\mathbf{x}_k) \right),$$

where  $0 \le \gamma \le \frac{2}{L}$ 

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PGM with exact line search

$$\gamma_{k} = \arg\min_{\gamma>0} f\left[ \mathsf{prox}_{\gamma b} \left( \mathsf{x}_{k} - \gamma \nabla a(\mathsf{x}_{k}) \right) \right]$$
$$\mathbf{x}_{k+1} = \mathsf{prox}_{\gamma b} \left( \mathsf{x}_{k} - \gamma_{k} \nabla a(\mathsf{x}_{k}) \right)$$
Form of bounds: Performance  $metric(\mathbf{x}_{k+1}) \leq t^*$  Performance  $metric(\mathbf{x}_k)$ Let  $\mathbf{g}_k$  denotes a (sub)gradient of f at  $\mathbf{x}_k$ , and

$$\rho_{\gamma} := \max\{\left|1 - \gamma \mu\right|, \left|1 - \gamma L\right|\}$$

Method	Step size $\gamma$	Metric	Rate t*	Known?
PGM	$\gamma \in (0, \frac{2}{L})$	$\left\ \mathbf{g}_{k}\right\ ^{2}$	$ ho_{\gamma}^2$	Taylor et al. (2018)

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Conclusion

Introduced an SDP hierarchy for bounding convergence rates via SOS certificates

 $\Rightarrow$  First level coincides with the PEP

#### Future work

- Other function classes
- Other algorithms
- Understanding when SOS certificates cannot be found

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### Future work

- Other function classes
- Other algorithms
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### Links

arXiv paper: http://arxiv.org/abs/1906.04648
github codes: https://github.com/sandratsy/SumsOfSquares

# Thank you!

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