

Canonical Estimation in a Rare Events Regime

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SILO (Oct 2011)

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- X_1, \dots, X_n are independent samples from p
- Law of large numbers:

$$\frac{1}{n} \sum_{i=1}^n X_i \xrightarrow{P} \mathbb{E}X$$

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Definition (Wagner, Viswanath and Kulkarni, IT-Trans 2011)

We say that $\{(\mathcal{A}_n, p_n)\}_{n \in \mathbb{N}}$ is a **rare-events source** if

$$\frac{\check{c}}{n} \leq p_n(a) \leq \frac{\hat{c}}{n}, \quad \forall a \in \mathcal{A}_n$$

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and $\exists P$ such that $P_n \Rightarrow P$. Note $\text{supp}(P) \subseteq \mathcal{C} := [\check{c}, \hat{c}]$.

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- Can we estimate all **reasonable** quantities in a universal manner?

Canonical Estimation Problems

Let $\{Y_n\}_{n \in \mathbb{N}}$ be a sequence of real-valued random variables such that

- There exists **continuous** $f_n(x)$ that converge to $f(x)$ **pointwise** on \mathcal{C}

$$\mathbb{E}[Y_n] = \int_{\mathcal{C}} f_n(x) dP_n(x)$$

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Definition

An estimator $\{\hat{Y}_n : \mathcal{A}_n^n \rightarrow \mathbb{R}\}_{n \in \mathbb{N}}$ is **consistent** if

$$\hat{Y}_n(X_{n,1}, \dots, X_{n,n}) \rightarrow \int_{\mathcal{C}} f(x) dP(x)$$

almost surely.

All previous examples were canonical

■ Probabilities

$$Y_n = \frac{1}{n} \log p_n^n(X_{n,1}, \dots, X_{n,n}) + \log n, \quad f(x) = \log x$$

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$$C = [\check{c}, \hat{c}], \quad f(x) = x^q, \quad \hat{c} = \lim_{q \rightarrow \infty} \left[\int_C x^q dP(x) \right]^{1/q}$$

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- Then by integrating against the correct limiting function f ,

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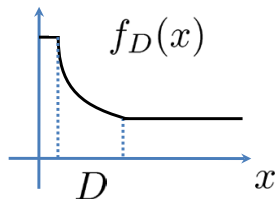
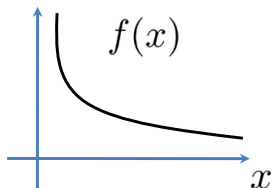
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- **Problems:**
 - Support \mathcal{C} isn't known
 - $f(x)$ doesn't have to be bounded everywhere
 - How to get the estimate $\hat{P}_n(x)$?

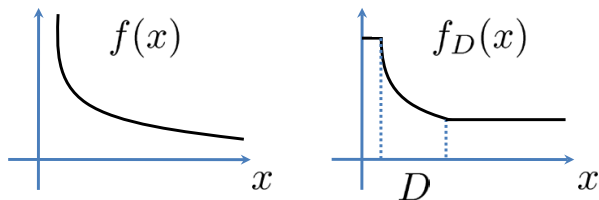
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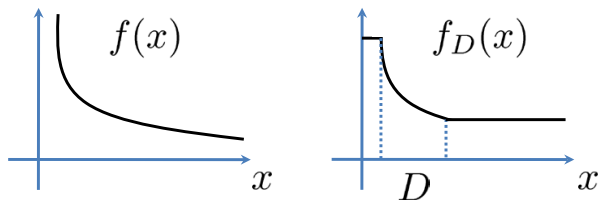
If $\mathcal{C} \subset \mathcal{D}$, then

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If \mathcal{C} is unknown, just **let \mathcal{D} grow gradually with n**

Estimate with rates

Recall the Wasserstein distance

$$d_W(P, Q) = \sup_{h \in \text{Lip}(1)} \left| \int_{\mathbb{R}^+} h dP - \int_{\mathbb{R}^+} h dQ \right|$$

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Lemma (Ohannessian-Tan-Dahleh)

If

$$\text{Lip}(f_{D_n}) d_W(\hat{P}_n, P) \rightarrow 0,$$

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How to estimate the shadow P_n ?

Pseudo-Empirical Measure

Good-Turing estimator:

- Denote the set of symbols that appear k times as $\mathcal{B}_{n,k} \subset \mathcal{A}_n$
- Denote their probabilities as

$$\gamma_{n,k} = p_n(\mathcal{B}_{n,k}) = \sum_{a \in \mathcal{B}_{n,k}} p_n(a)$$

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- E.g.: Probability of missing mass $\approx \phi_{n,0}$ [Badianu and Tong 2004]

Pseudo-Empirical Measure

Strong law of large numbers gives:

Lemma (WVK)

Let the P -Poisson mixture be

$$\lambda_k^P = \int_{\mathbb{R}^+} \frac{e^{-x} x^k}{k!} dP(x), \quad k = 0, 1, \dots$$

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Then, $\|\gamma_n - \lambda^P\|_1 \rightarrow 0$ and $\|\phi_n - \lambda^P\|_1 \rightarrow 0$ almost surely.

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Theorem (Ohannessian-Tan-Dahleh)

For “most natural” rare event sources, there exist an $s > 0$ such that

$$n^s \sup_{k \in \mathbb{N}} |F_{\phi_n}(k) - F_{\lambda^P}(k)| \rightarrow 0, \quad \text{a.s.}$$

(Kolmogorov-Smirnov convergence)

Estimation of $\hat{P}_n(x)$ via mixture distribution learning

Theorem

The (pseudo) *maximum-likelihood* estimator

$$\hat{P}_n^{ML} = \arg \min_Q D(\phi_n || Q)$$

is a valid construction, i.e., $\hat{P}_n^{ML} \Rightarrow P$ almost surely.

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Theorem

The *minimum distance* estimator

$$\hat{P}_n^{MD} = \arg \min_Q \sup_{k \in \mathbb{N}} |F_{\phi_n}(k) - F_{Poi(Q)}(k)|$$

is also valid. Furthermore, there exists $s > 0$ such that $n^s d_W(\hat{P}_n, P) \rightarrow 0$ almost surely (with some technical conditions).

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Lemma

With $D_n = o(n^s)$,

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Estimating support $\mathcal{C} = [\check{c}, \hat{c}]$

- Not quite canonical but close.
- Let Z be the weak limit of $Z_n := np_n(X_n)$ and let $P := Z_*(\mathbb{P})$. Then

$$\hat{c} = \operatorname{esssup}_{\omega} Z(\omega) = \lim_{q \rightarrow \infty} \left[\int_{\mathcal{C}} x^q dP(x) \right]^{1/q}$$

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Lemma

With $q_n = \frac{\log n}{\log \log n}$ and $D_n = o(n^{\frac{s}{2q_n}})$,

$$\hat{Y}_n := \left[\int_{\mathbb{R}^+} (x^{q_n})_{D_n} d\hat{P}_n(x) \right]^{1/q_n}$$

is consistent for estimating \hat{c}

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- **Future work:** Further analysis of convergence rates