

Minimax Optimal Fixed-Budget Best Arm Identification (BAI) in Linear Bandits.

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Setup: Arm set $[K] = \{1, \dots, K\}$ finite

Each $i \in [K]$ is assoc. to a vector $a(i) \in \mathbb{R}^d$

Arm vectors $\{a(1), a(2), \dots, a(K)\} \subset \mathbb{R}^d$
known

At each $t \in \mathbb{N}$, agent chooses $A_t \in [K]$, observes

$$X_t = \langle \theta^*, a(A_t) \rangle + \eta_t \quad (\text{Linear bandit})$$

$\theta^* \in \mathbb{R}^d$: unknown zero-mean 1-subG

\mathcal{E} : set of instances defined above $\{a(1), \dots, a(K), \theta^*\}$

Fixed-Budget # arm selections $T \in \mathbb{N}$ fixed

Online alg $\pi = (\pi_t)_{t=1}^T$ depends only on previous arm pulls rewards.

$\pi_t(a_t | A_1, X_1, A_2, X_2, \dots, A_{t-1}, X_{t-1})$

Goal: Design π that achieve

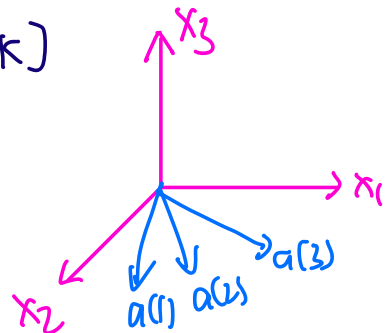
$$\min_{\pi} \Pr \left(i_{\text{out}} \neq \underbrace{\operatorname{argmax}_{j \in [K]} \langle \theta^*, a(j) \rangle}_{\text{unique best arm}} \right)$$

Observation: If $a(1), \dots, a(K)$ do not span \mathbb{R}^d , $\exists B \in \mathbb{R}^{d \times d}$ s.t.

$$a'(j) = B^T a(j), \quad j \in [K]$$

$$\{a'(j)\}_{j \in [K]} \text{ span } \mathbb{R}^d$$

B : calculated via a Gram-Schmidt.



OLS: $A_1, \dots, A_n \in [K]$ observe $X_1, \dots, X_n \in \mathbb{R}^d$

$a(A_1), \dots, a(A_n)$ span \mathbb{R}^d

$$\hat{\theta} = V^{-1} \sum_{t=1}^n a(A_t) X_t$$

$$V = \sum_{t=1}^n a(A_t) a(A_t)^T \text{ invertible}$$

Thm: If A_1, \dots, A_n are chosen deterministically, then $\forall b \in \mathbb{R}^d, \delta > 0$,

$$\Pr \left(\langle \hat{\theta} - \theta^*, b \rangle \geq \sqrt{2 \|b\|_{V^{-1}}^2 \log \left(\frac{1}{\delta} \right)} \right) \leq \delta$$

$$\|b\|_{V^{-1}}^2 = b^T V^{-1} b$$

Inspired by this

Aim: Find a prob. distn $\pi: \{a(i) : i \in [K]\} \rightarrow [0, 1]$ that

Min $\min_{\pi} g(\pi) = \max_{i \in [K]} \|a(i)\|_{V(\pi)^{-1}}^2$

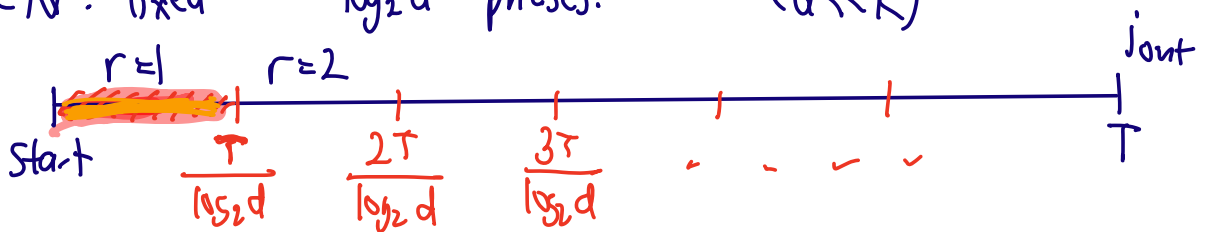
$V(\pi) = \sum_{i=1}^K \pi(a(i)) a(i) a(i)^T$

G-optimal design

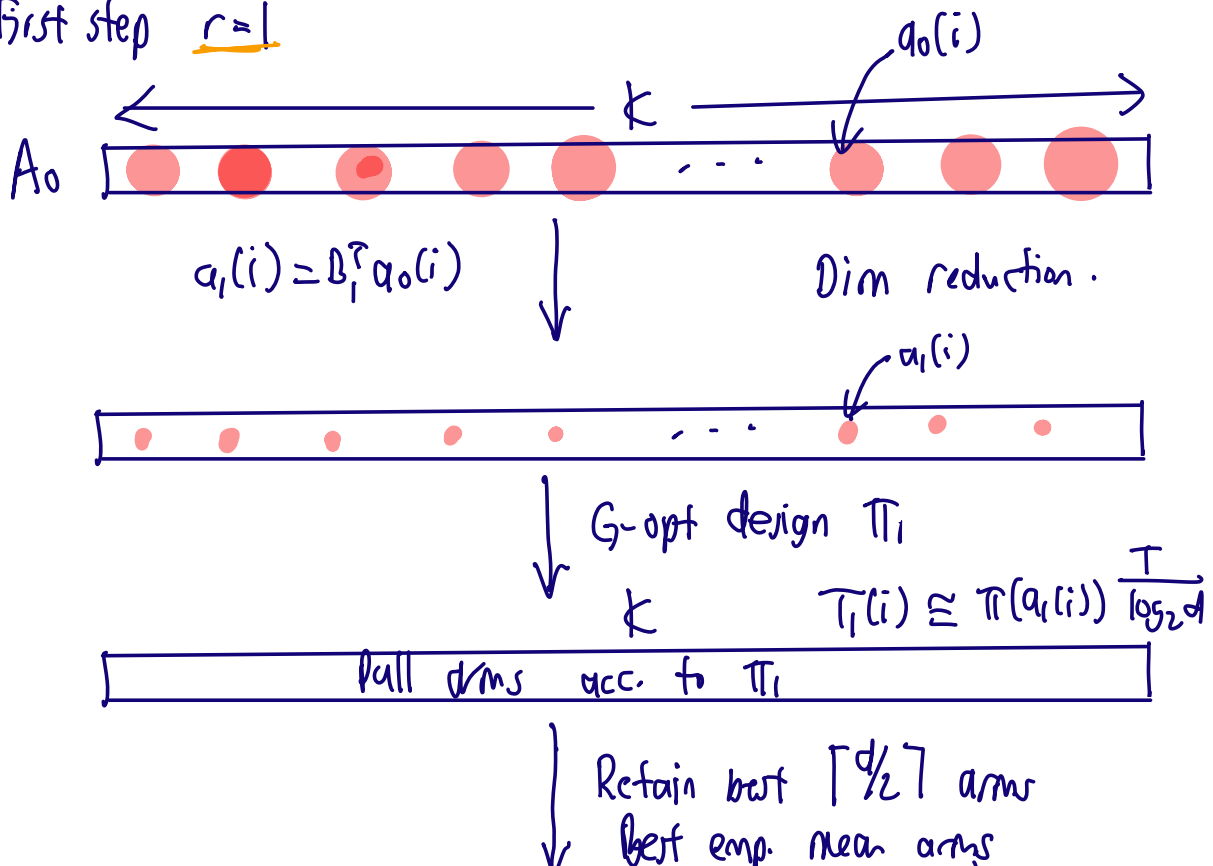
$g(\pi)$: related to confidence bds in the linear BAI prob.

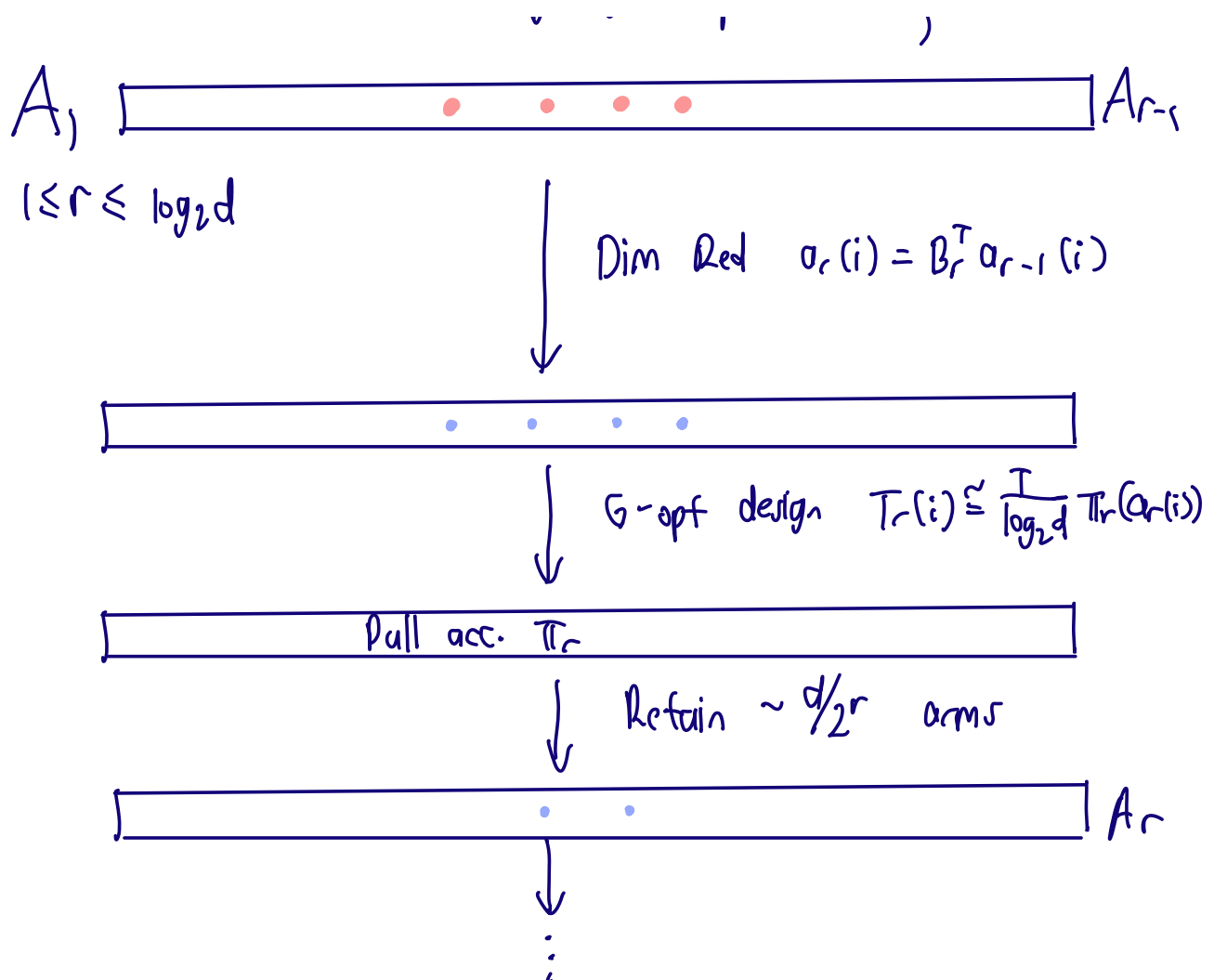
Algorithm: Optimal Design - Based Linear BAI (OD-LinBAI)

$T \in \mathbb{N}$: fixed $\log_2 d$ phases. ($d \ll K$)



First step $r=1$





Theoretical Guarantee for OD-LinBAI.

$$p(i) = \langle \theta^*, a(i) \rangle \quad p(1) > p(2) \geq \dots \geq p(K)$$

Define: $\Delta_i = p(1) - p(i), i \in [K] \setminus \{1\}$

$$\Delta_1 = \Delta_2, \quad \Delta_i \geq 0$$

Hardness Quantities $H_{1,\text{lin}} \stackrel{\Delta}{=} \sum_{i=1}^d \Delta_i^2, H_{2,\text{lin}} = \max_{i=1, \dots, d} \frac{i}{\Delta_i^2}.$

$$H_{2,\text{lin}} \leq H_{1,\text{lin}}$$

Thm: $\Pr(i_{\text{out}}^{\text{OD-LinBAI}} \neq 1) \lesssim \exp\left(-\frac{T}{32 \cdot H_{2,\text{lin}} \cdot \log_2 d}\right)$

Comparison with SOTA

Bayes Gap (Hoffman 2014) \Rightarrow not param-free

Peace (Katz-Samuel et al. NeurIPS 2020)

- i) not minimax opt.
- ii) not param-free

$$\sim \exp\left(-\Omega\left(\frac{T}{H_2 \log^2 d}\right)\right)$$

Linear Exploration (Alieva et al. ICML 2021)

$$\sim \exp\left(-\Omega\left(\frac{T}{H_2 \log_2 K}\right)\right)$$

GSE (Azizi et al. IJCAI 2022)

$$\sim \exp\left(-\Omega\left(-\frac{T \Delta_1^2}{d \log_2 K}\right)\right)$$

Lower Bd

$$\forall v \in \mathcal{E} \quad H_{\text{lin}} = H_{\text{lin}}(v)$$

$$\mathcal{E}(c) = \left\{ v \in \mathcal{E} : H_{\text{lin}}(v) \leq c \right\} \quad \text{Hardness bounded instances.}$$

Thm: For all T, c suff. large,

$$\inf_{\pi} \sup_{v \in \mathcal{E}(c)} \Pr(i_{\text{out}}^{\pi} \neq \underset{\substack{\uparrow \\ \text{best arm}}}{1}) \exp\left(\frac{2700T}{H_{\text{lin}}(c) \log_2 d}\right) \geq \frac{1}{6}$$

Rmk: Compare upper & lower bds.

OD-LinBAI: $\forall v \in \mathcal{E}$

$$\Pr(i_{\text{out}}^{\text{OD-LinBAI}} \neq 1) \leq \exp\left(-\Omega\left(\frac{T}{H_{2,\text{lin}} \cdot \log d}\right)\right)$$

all instances

Lower bd:

← specific instance

$$\Pr(i_{\text{out}}^{\pi} \neq 1) \geq \exp\left(-O\left(\frac{T}{H_{1,\text{lin}} \cdot \log d}\right)\right)$$

$H_{2,\text{lin}} \leq H_{1,\text{lin}} \Leftrightarrow \text{OD-LinBAI}$ is minimax optimal.

$\forall v \in \mathcal{E}$,

$$\Pr(i_{\text{out}}^{\text{OD-LinBAI}} \neq 1) \leq \exp\left(-\Omega\left(\frac{T}{H_{2,\text{lin}}(v) \cdot \log d}\right)\right)$$

$\forall \pi$ s.t. $\exists v \in \mathcal{E}(c)$,

$$\Pr(i_{\text{out}}^{\pi} \neq 1) \geq \exp\left(-O\left(\frac{T}{H_{1,\text{lin}}(v) \cdot \log_2 d}\right)\right)$$

For $\pi = \text{OD-LinBAI}$, $\exists v \in \mathcal{E}(c, \pi)$ s.t.

$$\Pr(i_{\text{out}}^{\pi} \neq 1) \geq \exp\left(-O\left(\frac{T}{H_{1,\text{lin}}(v) \cdot \log d}\right)\right)$$

This instance v must be s.t. $H_{1,\text{lin}}(v)$ is of the same order as $H_{2,\text{lin}}(v)$.

$$H_{1,\text{lin}}(v) = \Theta(H_{2,\text{lin}}(v)).$$

$$\forall \pi \quad \sup_{v \in \mathcal{E}(c)} \Pr(i_{\text{out}}^{\pi} \neq 1) \exp\left(\frac{2700T}{H_{1,\text{lin}}(v) \cdot \log_2 d}\right) \geq \frac{1}{6}$$

$\forall \pi \quad \exists v \in \mathcal{E}(c)$ s.t.

$$\Pr(i_{\text{out}}^{\pi} \neq 1) \exp\left(\frac{2700T}{H_{1,\text{lin}}(v) \cdot \log_2 d}\right) \geq \frac{1}{6}$$

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